# **Rhotrix-Decomposition**

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# Abstract

Matrix is decomposable if it can be expressed as a product of any two non singular lower and upper triangular matrices. A number of results on matrix decompositions are known in the literature. In this paper we introduce rhotrix decomposition which is a factorization of rhotrix into a product of rhotrices and also present some of its result.

**Keywords:** Rhotrix, Rhotrix multiplication, Rhotrix decompositions, Left and Right triangular rhotrices. AMS Subject Classifications [2010]: 15A15, 15A18

# 1.0 Introduction

The concept of rhotrices was first introduced by Ajibade [1] as an extension of the initiative on matrix-tertions and matrixnoitrets suggested by Atanassov and Shannon [2] and several works with modifications have presented on rhotrices. Ajibade [1] discussed the initial algebra and analysis on rhotrices and also set up some relationships between rhotrices and their hearts. The multiplication of rhotrices defined by Ajibade [1] was modified to similar multiplication of matrices proposed by Sani [3] defined as row column multiplication as follows: let A and B be any two rhotrices defined as

$$A = \begin{pmatrix} a_{11} & b_{11} \\ a_{21} & h(A) & a_{12} \\ a_{22} & a_{22} \end{pmatrix} \text{ and } B = \begin{pmatrix} b_{11} & b_{12} \\ b_{21} & h(B) & b_{12} \\ b_{22} & b_{22} \end{pmatrix} \text{ then } A \cdot B \text{ is defined as}$$
$$A \cdot B = \begin{pmatrix} a_{11}b_{11} + a_{12}b_{21} \\ a_{21}b_{11} + a_{22}b_{21} & h(A) \cdot h(B) & a_{11}b_{12} + a_{12}b_{22} \\ a_{21}b_{12} + a_{22}b_{22} \end{pmatrix}$$

This alternative multiplication method was then generalized for n-dimensional rhotrices by Sani [4]. The idea of the conversion of rhotrices to 'coupled matrices' was suggested by Sani in [5]. This idea was used to solve systems of  $n \times n$  and  $(n-1) \times (n-1)$  matrix problems simultaneously. This method was proved to be a linear transformation by Aminu [6]. The concept of vectors, one-sided system of equations and eigenvector-engenvalue problem in rhotrices were introduced by Aminu [7]. A necessary and sufficient condition for the solvability of one sided system of rhotrix was also presented in [8]. It was shown in the paper how a solution can be found provided the system is solvable. Rhotrix vector spaces and their

properties were presented by Aminu [9]. Cayley–Hamilton theorem in rhotrix, determinant method for solving system of equation in rhotrices and minimal polynomial of a rhotrix were discussed in [10,11,12] respectively.

In this article, we will discuss and present the necessary and sufficient conditions for the solvability of such systems. If this system is solvable, we show how a solution can be found.

Decomposition of some special rhotrix 'Vandamonde Rhotrix' was discussed in Kumar and Sharma [13].

Our aim in this chapter is to present rhotrix decomposition, a concept which to the best of our knowledge has not been studied before.

Rhotrix has applications several applications in coding theory, cryptography, combinatorial design graph theory, see Kumar and Sharma [14, 15].

### **Definition 1.1**

An *n*-dimensional rhotrix is represented and given as follows:

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$$\mathcal{R}_{n} = \langle a_{ij}, c_{lk} \rangle = \begin{pmatrix} a_{21} & c_{11} & a_{12} & c_{12} & a_{13} & c_{12} & c_{12} & a_{13} & c_{12} & c_{12} & a_{13} & c_{12} & c_{12} & c_{12} & a_{13} & c_{12} & c_{$$

Where t = (n + 1)/2. The multiplication method of rhotrices is a similar as that of matrices.

**Definition 1.2**The heart or centre of a rhotrix is the element that lies at the middle of a given rhotrix thereby dividing it into two equal parts. It lies at the position  $\frac{1}{2} \left[ \frac{1}{2} (n^2 + 1) + 1 \right]$ , Where *n* is the dimension of the rhotrix.

**Definition 1.3**A system of linear equations in rhotrix form is given by:

 $\mathcal{R}_n \cdot \langle x^{nj} \rangle = \langle b^{nj} \rangle$ 

Where  $\mathcal{R}_n = \langle a_{ij}, c_{lk} \rangle$  for: *i*, *j* = 1,2, ... *t*, and *k*, *l* = 1, 2, ..., *t* - 1.

**Definition 1.4**The rhotrix defined in Definition 1.1 above is said to be **left triangular rhotrix** or simply**left rhotrix** if  $a_{ij} = 0$  for  $2 < i \le t$  and  $1 \le j \le t$  for t = 1, 2, ..., n. We denote  $\mathbf{R}^l_n$  to represent *n* dimensional left rhotrix. Examples of 3

and 4 dimensional left rhotrix is given as  $\mathbf{R}^{l}_{3} = \begin{pmatrix} a_{11} \\ a_{21} \\ a_{22} \end{pmatrix}$  or  $\mathbf{R}^{l}_{5} = \begin{pmatrix} a_{11} \\ a_{21} \\ a_{22} \\ a_{11} \\ a_{12} \\ 0 \end{pmatrix}$  respectively.

And n dimensional left rhotrix is

$$\mathcal{R}_{n}^{\ l} = \begin{pmatrix} l_{21} & c^{l}{}_{11} & 0 \\ l_{31} & c^{l}{}_{21} & l_{22} & 0 & 0 \\ l_{t1} & \vdots & \vdots & \vdots & \vdots & \vdots & 0 \\ \vdots & 0 \\ l_{t1-2} & l_{(t-1)(t-2)} & l_{(t-1)(t-1)} & 0 & 0 & \cdots \\ l_{tt-1} & c^{l}{}_{t-1t-1} & 0 \\ l_{tt} & l_{tt} &$$

Where  $c^{l}$ , is an element in the **left** rhotrix.

**Definition 1.5** The rhotrix defined in Definition 1.1 above is said to be **right triangular rhotrix** or simply **right rhotrix** if  $a_{ij} = 0$  for  $2 < j \le t$  and  $1 \le i \le t$  for t = 1, 2, ..., n. Here we denote  $\mathbf{R}^r_n$  to represent *n* dimensional right rhotrix. An

example of 3 dimensional right rhotrix is given as:

$$\boldsymbol{R}_3 = \begin{pmatrix} a_{11} \\ 0 & r & a_{12} \\ a_{22} \end{pmatrix}$$

An *n* dimensional right rhotrix is

$$\mathcal{R}_{n}{}^{r} = \begin{pmatrix} & 0 & c^{r}{}_{11} & r_{12} \\ & 0 & c^{r}{}_{12} & c^{r}{}_{12} & r_{13} \\ & \cdots & 0 & \cdots & \cdots & \cdots & \cdots \\ 0 & \vdots & \vdots & \vdots & \vdots & \vdots & r_{1t} \\ & \cdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots \\ & 0 & 0 & r_{(t-1)(t-1)} & r_{(t-2)(t-1)} & r_{(t-2)t} & \\ & 0 & c^{r}{}_{t-1t-1} & r_{t-1t} \\ & & r_{tt} & & & & \\ \end{pmatrix}$$

Where  $c^r$ , is an element in the **right** rhotrix.

**Definition 1.6**An *n* dimensional rhotrix  $\mathcal{R}_n$ , is said to have an  $\mathbf{R}^r_n \cdot \mathbf{R}^l_n$  or simply  $\mathbf{R}\mathbf{L}$  rhotrix decomposition if there exists  $\mathbf{R}^r_n$  and  $\mathbf{R}^l_n$  rhotriceswhere  $\mathbf{L}$  and  $\mathbf{R}$  are left and Right rhotrices respectively. This is defined as  $\mathcal{R}_n = \mathbf{R}^r_n \cdot \mathbf{R}^l_n$  where

$$\boldsymbol{R}^{l}{}_{n}\boldsymbol{R}^{r}{}_{n} = \begin{pmatrix} \begin{matrix} l_{21} & c^{l}{}_{11} & 0 & & & r_{11} \\ l_{21} & c^{l}{}_{11} & 0 & & & 0 & c^{r}{}_{11} & r_{12} \\ \\ l_{21} & c^{l}{}_{21} & l_{22} & 0 & 0 & & \\ \\ l_{t1} & \vdots & \vdots & \vdots & \vdots & \vdots & 0 \\ \\ l_{t1} & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & 0 \\ \\ l_{t1-2} & l_{(t-1)(t-2)} & l_{(t-1)(t-1)} & 0 & 0 & \\ \\ l_{tt-1} & c^{l}{}_{t-1t-1} & 0 & & & 0 & c^{r}{}_{(t-1)(t-1)} & r_{(t-2)(t-1)} & r_{(t-2)t} & \\ \\ l_{tt} & & & & l_{tt} & & \\ \end{matrix} \right) \cdot \begin{pmatrix} \dots & \dots & n & n \\ 0 & \vdots & \vdots & \vdots & \vdots & r_{1t} \\ \dots & \vdots & \vdots & \vdots & \vdots & \vdots & r_{1t} \\ \\ \dots & 0 & 0 & r_{(t-1)(t-1)} & r_{(t-2)(t-1)} & r_{(t-2)t} & \\ \\ 0 & 0 & c^{r}{}_{t-1t-1} & r_{t-1t} & \\ \\ r_{tt} & & & r_{tt} & \\ \end{pmatrix}$$

$$\mathbf{R}^{r}_{n} \cdot \mathbf{R}^{l}_{n}$$

$$= \begin{pmatrix} l_{11}r_{11} & l_{11}r_{12} & l_{11}r_{13} & l_{11}r_{14} &$$

#### 2.0 LR–Rhotrix Decomposition

Given a System of linear equations:  $\mathcal{R}_n X = b$ , let  $\mathcal{R}_n = LR$  where  $R_n^l$  and  $R_n^r$  are Left and Right triangular rhotrices. Let denotes L and R as follows: here we give more consideration for n = 3, which is the order of the rhotrices. Note that, we will be using  $\mathcal{R}_n^l = L$  and  $\mathcal{R}_n^r = R$ , throughout the work, represent, left and right rhotrices respectively.

$$\boldsymbol{L} = \begin{pmatrix} l_{11} \\ l_{21} \\ l_{22} \end{pmatrix} \text{ and } \boldsymbol{R} = \begin{pmatrix} 0 & r_{11} \\ r_{12} \\ r_{22} \end{pmatrix}$$
Then,
$$(2.1)$$

$$\boldsymbol{L}\boldsymbol{R} = \begin{pmatrix} l_{11} \\ l_{21} \\ l_{22} \end{pmatrix} \cdot \begin{pmatrix} \sigma \\ r \\ r_{22} \end{pmatrix} = \begin{pmatrix} l_{11}r_{11} \\ l_{21}r_{11} \\ lr \\ l_{21}r_{12} + l_{22}r_{22} \end{pmatrix} = \boldsymbol{\mathcal{R}}_{3}$$
(2.2)

Crout's and Doolittle's methods were first introduced on matrices and have received a number of attentions; the reader is referred to Bunch [16] for more information. Here we need to define Crout's and Doolittle's Methods using rhotrices.

#### **Crout's Algorithm for Rhotrices:** 2.1

Using (2.2) above we have:  $r_{ii} = r = 1$ , i = 1, 2. Then

$$LR = \begin{pmatrix} l_{11} \\ l_{21} \\ l_{21}r_{12} + l_{22} \end{pmatrix} = \mathcal{R}_3$$
(2.3)

In **Crout's Methodfor rhotrix**, we use:  $r_{ii} = r = 1$ , where r is the centre of the right rhotrix

#### 2.2 **Doolittle's Algorithm for Rhotrices:**

using (2.2) as above where  $l_{ii} = l = 1$ , then for i = 1, 2.

$$LR = \begin{pmatrix} l_{21}r_{11} & r & r_{12} \\ l_{21}r_{12}+r_{22} \end{pmatrix} = \mathcal{R}_3$$
(2.4)

Similarly, as in Crout Method for rhotrices, **Doolittle's Method for rhotrix** we use:  $l_{ii} = l = 1$ , where l is the centre of the left rhotrix.

For example, let's solve a system below using any of the above algorithms.

 $2x_1 - 3x_2 = 1$ A system

$$x_1 + x_2 = 3$$

Can be solve using any of these algorithms, but here in particular, we are going to deploy the first algorithm, which is Crout's Method.

The system in rhotrix form is

$$\begin{pmatrix} 2 \\ 1 \\ 4 \\ 1 \end{pmatrix} \cdot \begin{pmatrix} x_1 \\ x_2 \\ 0 \\ 0 \end{pmatrix} = \begin{pmatrix} 1 \\ 3 \\ 0 \\ 0 \end{pmatrix}$$
(2.5)

By choosing an arbitrary Heart, here choose a heart (4). This can be represented in an equivalent form as:  $\mathcal{R}_n \langle x^{nj} \rangle = \langle b^{nj} \rangle$ 

For j = 1, 2. n = 3We have:

 $\mathcal{R}_{3}\langle x^{31}\rangle = \langle b^{31}\rangle$ 

Where  $\langle x^{31} \rangle$  and  $\langle b^{31} \rangle$  are column vector rhotrices. Re-writing the rhotrix system in the form

 $\mathcal{R}_3 = LR$ , where L and R are Left and Right triangular rhotrices.

Now, using Crout's algorithms to rhotrices:  $r_{ii} = r = 1$ . Using (2.3) we have:

$$LR = \begin{pmatrix} l_{11} \\ l_{21} & l & l_{11}r_{12} \\ l_{21}r_{12} + l_{22} \end{pmatrix} = \begin{pmatrix} 2 \\ 1 & 4 \\ 1 \end{pmatrix} = \mathcal{R}_3$$

So we get:  $l_{11} = 2$ ,  $l_{21} = 1$ , l = 4,  $l_{11}r_{12} = -3$  which implies  $r_{12} = -\frac{3}{2}$ ,  $l_{21}r_{12} + l_{22} = 1$  or  $(1)\left(-\frac{3}{2}\right) + l_{22} = 1 \Rightarrow l_{22} = \frac{5}{2}$ . With  $r_{11} = r_{22} = 1$ . So we have:  $L = \left( l_{21} \quad l \quad 0 \\ l_{22} \quad l \quad 0 \\ l_{23} \quad l_{23} \quad l_{23} \\ l_{23} \quad l_{23$ 

We solve the system by **back substitution** by first solving the system below: So that:  $LR \cdot \mathcal{R}_3 \langle x^{31} \rangle = \langle b^{31} \rangle$ 

Let 
$$R \cdot \langle x^{31} \rangle = \langle y^{31} \rangle$$
 then  $L \cdot \langle y^{31} \rangle = \langle b^{31} \rangle$  now we have  

$$\begin{pmatrix} 2 \\ 1 & 4 \\ 5 \\ 2 \end{pmatrix} \cdot \begin{pmatrix} y_1 \\ y_2 & 0 \\ 0 \end{pmatrix} = \begin{pmatrix} 3 & 0 \\ 3 & 0 \\ 0 \end{pmatrix}$$

The matrix equivalent form is  $\begin{pmatrix} 2 & 0 \\ 1 & \frac{5}{2} \end{pmatrix} \begin{pmatrix} y_1 \\ y_2 \end{pmatrix} = \begin{pmatrix} 1 \\ 3 \end{pmatrix}$ ,

 $\Rightarrow 2y_1 = 1, \text{ or } y_1 = \frac{1}{2} \text{ and } y_1 + \frac{5}{2}y_2 = 3, \frac{5}{2}y_2 = 3 - \frac{1}{2} = \frac{5}{2} \text{ therefore } y_1 = \frac{1}{2} \text{ and } y_1 = 1.$ Now, **forward substitution:**solving for  $\langle x^{31} \rangle$  in the system :  $R \cdot \langle x^{31} \rangle = \langle y^{31} \rangle$ 

$$\begin{pmatrix} 1 & 1 & -\frac{3}{2} \\ 0 & 1 & -\frac{3}{2} \end{pmatrix} \cdot \begin{pmatrix} x_2 & 0 & 0 \\ 0 & 0 \end{pmatrix} = \begin{pmatrix} y_2 & 0 & 0 \\ 0 & 0 \end{pmatrix} = \begin{pmatrix} \frac{1}{2} & 0 \\ 1 & 0 & 0 \\ 0 & 0 \end{pmatrix}$$
  
This implies that  $\begin{pmatrix} 0 & 1 & -\frac{3}{2} \\ 1 & 0 & 0 \end{pmatrix} \cdot \begin{pmatrix} x_2 & 0 & 0 \\ 0 & 0 \end{pmatrix} = \begin{pmatrix} \frac{1}{2} & 0 \\ 1 & 0 & 0 \\ 0 & 0 \end{pmatrix}$ 

Converting to its matrix equivalent form: we have  $x_2 = 1$  and  $x_1 - \frac{3}{2}x_2 = \frac{1}{2}$  and  $x_2 = \frac{1}{2}$  and  $x_1 - \frac{3}{2}x_2 = \frac{1}{2}$  and  $x_2 = \frac{1}{2}$  and  $x_3 = \frac{1}{2}$ .

 $x_1 = \frac{1}{2} + \frac{3}{2} = \frac{4}{2} = 2.$  So  $x_1 = 2$  and  $x_2 = 1.$ 

Hence we clearly check the system for  $x_1 = 2$  and  $x_2 = 1$  the truth is clear.

# Theorem 2.1

A necessary and sufficient condition for an *n*dimensional rhotrix,  $\mathcal{R}_n$ , to be  $\mathbf{R}^r_n \cdot \mathbf{R}^l_n$  or simply  $\mathbf{L}\mathbf{R}$  – rhotrix decomposition is that  $\mathcal{R}_n$  is invertible.

# **Proof:**

Let  $\mathcal{R}_n = \langle a_{ij}, c_{lk} \rangle$  for  $1 \le i \le n, 1 \le j \le n$  and  $1 \le k \le n$ , suppose  $\mathcal{R}_n$  admits *LR* decomposition, we need to show that *det* ( $\mathcal{R}_n$ ) exists.

Since  $\mathcal{R}_n$  can be express as follows  $\mathcal{R}_n = LR$ , where L and R are Left and Right triangular rhotrices. Now, the  $det(L) \neq 0$  and  $det(R) \neq 0$  hence the determinant  $det(LR) \neq 0$  which implies that  $det(\mathcal{R}_n) \neq 0$ . Thus  $\mathcal{R}_n$  is invertible. Conversely, if  $\mathcal{R}_n$  is invertible its determinant exists. That is

$$det (\mathcal{R}_n) \neq 0 \Rightarrow det (LR) \neq 0$$

Now, since  $det(LR) \neq 0$  then  $det(L) \neq 0$  and  $det(R) \neq 0$ , thus  $\mathcal{R}_n$  can be express in the form  $\mathcal{R}_n = LR$ . Hence LR decompositions exist when  $\mathcal{R}_n$  is invertible.

**Definition 2.1**An *LDR*-rhotrix decomposition is of the form:  $\mathcal{R}_n = LDR$ , where *D* is a <u>diagonal rhotrix</u> and *L* and *U* are *unit* triangular rhotrices, meaning that all the entries on the diagonals of *L* and *U* are one. It is equivalent to *LR* rhotrix decomposition.

# Theorem 2.2

If an invertible rhotrix has an LDR rhotrix factorization, then its unique, in that case, the LR factorization is also unique if we require the diagonal of L or R consists of ones.

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#### **Proof:**

By using any of the algorithms, Crout or Doolittle Methods we can arrive at the proof.

Now, using Crout's method where  $r_{ii} = r = 1$  and *D* a diagonal rhotrix

Suppose  $\mathcal{R}_n$  has LDR rhotrix factorization, then it can be expressed in the form:

 $\mathcal{R}_n = RDL = \langle x^{ii} \rangle \cdot \langle x^{nn} \rangle \cdot \langle x^{jj} \rangle$  where *D* is a diagonal rhotrix. Since  $\mathcal{R}_n$  is invertible it follows from theorem 1.1 that  $\mathcal{R}_n = RDL$  exists.

And since it exists then RDL = LDR hence the factorization is unique.

It then follows from RDL = LDR

Pre-multiplying through by  $D^{-1}(RDL) = D^{-1}(LDR)$  where

$$\hat{R}D^{-1}DL = LD^{-1}DR$$

$$RL = LR$$

Hence **RL** factorization is unique provided that **LDR** factorization exists.

### Theorem 2.3

If  $\mathcal{R}_n$  is a non-singular rhotrix, then there exists a permutation rhotrix *P* so that  $P\mathcal{R}_n$  has an *LR*- decomposition. i.e.  $P\mathcal{R}_n = LR$ .

**Proof:** 

Or

Since  $\mathcal{R}_n$  is invertible being non singular, it follows from theorem 1.2. note we use *P* only when  $\mathcal{R}_n$  requires row interchanges to row echelon form, and  $P\mathcal{R}_n$  requires no row interchanges. So we have  $P\mathcal{R}_n$  can be express as  $P\mathcal{R}_n = LR$ . Hence the proof.

**Definition 2.2**An n dimensional rhotrix which arises by a finite number of row interchange is called a permutation rhotrix. It is usually denotes as P.

**Theorem 2.4**The system  $\mathcal{R}_n X = b$  is equivalent to the system  $P\mathcal{R}_n X = Pb$ , where  $P\mathcal{R}_n$  has *LR* decomposition and *P* is a permutation rhotrix.

### Proof:

We need to show that the system  $\mathcal{R}_n X = b$  is equivalent to  $P\mathcal{R}_n X = Pb$ . It suffices to verify that if  $P\mathcal{R}_n X = Pb$ , P interchanges the rows of the rhotrix  $\mathcal{R}_n$ ,

Now, if  $P\mathcal{R}_n X = Pb$  is given: pre-multiply both sides with  $P^{-1}$  we get:

$$P^{-1}P\mathcal{R}_n X = P^{-1}Pb$$

 $\Rightarrow IP\mathcal{R}_n X = Ib \text{ or } \mathcal{R}_n X = b$ 

Hence  $\mathcal{R}_n X = b$  is equivalent to: $P\mathcal{R}_n X = Pb$ .

Now, if  $\mathcal{R}_n X = b$  is given, to interchange the rows of  $\mathcal{R}_n$  into row echelon form we multiply through by the permutation rhotrix *P*, such that  $P\mathcal{R}_n X = Pb$ .

Hence  $\mathcal{R}_n X = b \Leftrightarrow P \mathcal{R}_n X = P b$ 

# Theorem 2.5

Let  $\mathcal{R}_n X = b$  be a rhotrix system of linear equation where  $\mathcal{R}_n$  has zero heart. Using LR decomposition of rhotrices, by applying either Crout or Doolittle methods the heart of either L or R must be zero.

# Proof:

Let l and r be the centres (hearts) of L and Rleft and right triangular rhotrices respectively.

If  $\mathcal{R}_n = LR$ , where the initial system in matrix form is transformed to rhotrix form has zero heart, so the product of the their hearts gives the heart of  $\mathcal{R}_n$ . If the heart of  $\mathcal{R}_n$  is 0 then we have:  $r \cdot l = 0$  which implies that either, r = 0 or l = 0, Hence the proof.

# **Definition 2.3**

The total number of floating point operations  $(\times, \div, +, -)$  determine the cost of computations involved in solving problems. It's usually denoted as  $O(n^k)$ , for natural number  $k \ge 1$  and n is the dimension of the rhotrix.

# Theorem 2.6

An *LR* rhotrix decomposition requires  $O(n^3)$  floating point operations.

#### Proof:

In order to compute the total number of operations we will need the following identities:

$$\sum_{i=1}^{n} 1 = n, \quad \sum_{i=1}^{n} i = \frac{n}{2}(n+1), \quad \sum_{i=1}^{n} i^2 = \frac{n}{6}(n+1)(2n+1)$$
 which can be proved using induction

There are (n + 1 - i) and (n - i + 2) multiplication(×) and division(÷) operations respectively.

Therefore the total number of multiplication( $\times$ ) and division( $\div$ ) operations is:

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$$\sum_{i=1}^{n-1} (n+1-i)(n-i+2) = (n^2+3n+2) \sum_{i=1}^{n-1} (1-(2n+3)) \sum_{i=1}^{n-1} (1+i)(n-i+2) = (n^2+3n+2) - (2n+3) \frac{n}{2}(n+1) + \frac{n}{6}(n+1)(2n+1) = \frac{1}{3}n^3 + \frac{1}{2}n^2 + \frac{2}{3}n \approx \frac{1}{3}n^3 + \frac{1}{3}n^$$

Similarly,

There are (n - i) and (n - i + 1) addition(+) and subtraction(-) operations respectively.

Therefore the total number of addition(+) and subtraction(-) operations is:

$$\sum_{i=1}^{n-1} (n-i)(n-i+1) = (n^2+n) \sum_{i=1}^{n-1} 1 - 3 \sum_{i=1}^{n-1} i + \sum_{i=1}^{n-1} i^2 = \frac{1}{3}n^3 + \frac{1}{2}n^2 - \frac{5}{6}n \approx \frac{1}{3}n^3.$$

The approximate total solution cost is:  $=\frac{1}{3}n^3 + \frac{1}{2}n^2 + \frac{2}{3}n + \frac{1}{3}n^3 + \frac{1}{2}n^2 - \frac{5}{6}n = \frac{2}{3}n^3 + n^2 - \frac{1}{6}n \approx \frac{2}{3}n^3$ .

So,  $O(n^3) \approx \frac{2}{3}n^3$ .

# Note:

- 1. If two rhotrices of order *n* can be multiplied in  $\mathcal{R}(n)$ , where  $\mathcal{R}(n) \ge n^a$  for some a > 2, then the *LR* decompositions the rhotrix decompositions can be computed in  $O(\mathcal{R}(n))$ .
- 2. Computing the *LR* rhotrix decomposition of rhotrices using either of these algorithms requires  $O(n^3) = \frac{2n^3}{3}$  floating point operations, ignoring the lower order terms.

### 3.0 Conclusion

In this paper we havediscussed the concept of rhotrix decomposition and its properties.

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